



# Availability of Optical Network with Self-Healing Nodes Designed by Architecture on Demand

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# RELIABILITY AVAILABILITY SECURITY



#### What kind of quality do we expect from a system?





to work correctly or to be in idle secure state in required time



#### What kind of quality do you expect from a system? (2)



to work correctly or to be in idle secure state in required time





- System reliability *R*(*t*) is the probability that the system works correctly in the period of time *t* under defined environmental conditions.
- System availability A(t) is the probability that the system works correctly at the time point t.
- System security (safety) *S*(*t*) is the probability that the system works correctly or does not work at all in the period of time *t* under defined environmental conditions.





- Large volume of traffic is carried over transmission networks
- Loss of service means loss of income:
  - Direct losses, service cannot be charged
  - Indirect losses
    - Penalties "SLA not fulfilled"
    - Loss of clients

• **SLA** – Service Level Agreement



Availability – proper quality

measure in communications?



Availability  $\rightarrow$  Is it proper quality measure for **communication** systems and networks?

The most important question: What are the **consequences** of an **outage** of a communication system or network?

 $\rightarrow$  loss of money for operator and some troubles for subscribers  $\rightarrow$  **but not a catastrophic event**.

- **Low** availability  $\rightarrow$  high traffic losses
- High availability → low traffic losses, but high development, manufacturing and spare capacity costs.
- Optimal availability → compromise of quality and cost

For special communication services  $\rightarrow$  other requirements could be specified (i.e. real time remote control (*R*) or money transfer (*S*).



An example: Availability of

optical transport system?



**Example:** 1 optical cable/24 fibres/32  $\lambda$  per fibre, 60% utilization Single wavelength ( $\lambda$ ) (40 Gbit/s): SDH transport network 1×STM-256 = 256 STM-1 = 256×63  $E_1$  =

=  $256 \times 63 \times 30 E_0$  = 483,840 ch/ $\lambda$  ( $E_0$ =64 kbit/s)

Assuming channel ( $E_0$ ) interconnection price:  $C_{ch} = 0.01 \ C/ch/min$ 1 minute of a wavelength ( $\lambda$ ) unavailability (*LOR*\*per *min* per  $\lambda$ )

 $LOR_{min/\lambda} = 483,840 \times 0.01 \approx 5,000 \text{ C/min/}\lambda$ 

For entire cable: ×24×32×0.6 = 2,229,534.72 ≈

*LOR*<sub>min/cable</sub> ≈ 2 M€/min/cable

Availability of A=0.99999 (unavailability  $U=10^{-5}$ ) amounts  $MDT^{**}=5.256 \text{ min/year}$  (service unavailability per year)  $LOR_{year/cable} = 11,718,434.49 \in ~ 12 \text{ M} \text{C/year/cable}$ 

**LOR**\* – Loss of Revenue

*MDT*\*\* – Mean Down Time, (*MDT* = *U*×**525,600 min/year**)





# RELIABILITY & AVAILABILITY OF COMMUNICATION NETWORK



Definitions of communication network availability



A<sub>s,t</sub> s,t-availability the most common measure: "communication network availability" (non-linear measure)

*s,t*-availability of communication network is the worst node pair (s, t) availability among all pairs of nodes.

- $A_g$  **g**-availability probability that all *n* nodes in a network are connected (k=n)
- $A_{i,j}$  *i,j*-availability node pair (*i*, *j*) availability (**basic measure**)
  - average availability (linear measure)
- $A_k$ k-availabilityprobability that k nodes in a network<br/>are connected (if  $k=2 \rightarrow s,t$ -availability)<br/>(the most general measure)

 $A_{av} av$ -availability







$$S_i = \frac{dR}{dR_i} = \frac{\Delta R}{\Delta R_i} = \frac{R(R_i) - R(0)}{R_i} = \frac{R(1) - R(R_i)}{1 - R_i} = \frac{R(1) - R(0)}{1 - 0} = R(1) - R(0)$$

# **Protection & restoration scenarios**







# Case study: Availability of Optical Network with Self-Healing Nodes Designed by Architecture on Demand Can be Architecture on Demand Nodes Used to Improve Availability of Optical Networks?

Joint project of University of Zagreb and University of Bristol





Hard-wired (HW) optical node architectures

- Low level of architectural flexibility and scalability.
- Unnecessary components in the chain 
   lower availability
- Recovery after any node component failure is possible only on the network level
  - there is a need for number of protocol steps
  - switchover time in the range of seconds.





#### Usage of **additional switching elements**

(e.g. mirrors in 3DMEMS or collimator arrays in piezo electric switch)

#### **Disadventage:**

- Additional switches > higher node cost > lower availability
- **Adventages:**
- arbitrary interconnection of optical components
- additional functionality
- different switching granularities,

*fibre, waveband, wavelength* and *sub-wavelength* (time)













**Self-healing** after a failure of AoD node component is based on switchover to the **node redundancy**:

- intentionally added for survivability enhancement
- created by releasing unused components
  - ✤ At lower traffic load in a node → surplus of components are released and used as redundant
  - ✤ By grouping of wavelength or waveband paths into fibreswitched paths ⇒ node components are released and used as redundant.

#### Basic question: If the redundancy + self-healing ⇒ higher availability?



### Architecture on Demand (AoD) experiment



highperformancenetworks group

- Backplane provides connectivity and reconfiguration
- Modules provide required functionality, e.g. amplification, spectral demultiplexing, subwavelength switching, etc.





#### **Hard-wired nodes**







#### Node with Architecture on Demand







- All redundant components can be used on-the-fly for self-healing
  - provide backup resources for paths on the optical node level - switchover time at node level > 10 milisecond range (when 3D MEMS is applied)
  - no necessity for path rerouting on the optical network level - path recovery at network level switchover time in second range.
- Only necessary components are traversed by a path
   less components > higher path and network
   availability



#### Main trade-offs at AoD



Additional switches

VS.

Redundancy & Self-healing

- Lower availability & High AoD cost 
   Higher operator's revenue losses
- Higher power consumption

- Higher availability
   Lower operator's
   revenue losses
- Lower no. of active components
   Lower power consumption



#### Experimental setup –AOD ROADM with additional redundancy





- Restoration time ≈ 20 ms ✓
- Redundancy for any type of component without disturbing the existing connections  $\checkmark$
- Usage of expensive optical backplane (3D-MEMS) ×



### **Self-healing procedure**





- **1** Initial cross-connection (within 1<sup>st</sup> SP) from input 1 to output 2 at node **B**
- 2 Component WSS#1 fails
- 3 Self-healing: New cross-connections can use idle (redundant) WSS#2
- ④ If self-healing at B is not possible ➡ backup path (2<sup>nd</sup> SP) = protection path





- *HWnet* and *AoDnet* are compared using availability measures:
- *s,t-availability* (*A<sub>s,t</sub>*) minimal value of all *i,javailabilities* (*A<sub>i,j</sub>*) among all node pairs - represents the worst end-to-end connection in a network.
- *g-availability* (*A<sub>g</sub>*) is the probability that all end-to-end connections are in working state.
- Mean down time ( $MDT_g$ ) is the time, in minutes per year, when at least one end-to-end connection is broken:  $MDT_g = (1-A_g) \times 525,600$  minutes/year.





- 1<sup>st</sup> shortest and 2<sup>nd</sup> shortest paths are found for each source and destination node pair of connection requests.
- Each node is reconfigured for all paths which fulfil the requirements for fibre switching (FS). A SPL and one WSS are released and placed in the node spare bank.
- For reconfigured network, Monte Carlo failure/repair simulation is carried out.

## **Enforced fiber switching (EFS)**



- Fibre switching can be enforced via careful routing on the network level
- Main idea: Allow for fibre switching in nodes by rerouting the extra lightpaths

$$fs\_ratio \text{ (node } n, \text{ input } = \frac{\# \text{lightpaths which use port } i \text{ and port } j}{\# \text{lightpaths which use either port } i \text{ or port } j}$$
Must be rerouted to allow for fibre switching

$$fs\_ratio = \frac{2}{1} = 2$$

Can be fibre-switched between ports (*i*, *j*) in node n

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An example network – assumptions



- **German topology** (11 nodes and 34 links).
- Network protection scheme: **1+1**.
- The traffic load in the network is assumed to be **static**.
- **Populations & distance** traffic calculus for each node pair.
- Traffic requirements are fulfilled by multiples of **10 Gbit/s** lightpaths.
- No. of *SPLs* or *WSSs* = node degree.
- Each optical end-to-end connection 
   pair of paths:
  - first shortest path (1<sup>st</sup> SP) working path and
  - second shortest path (2<sup>nd</sup> SP) backup path (independent of the working path).



![](_page_30_Picture_1.jpeg)

![](_page_30_Figure_2.jpeg)

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31

![](_page_31_Picture_1.jpeg)

![](_page_31_Picture_3.jpeg)

- Monte Carlo failure/repair simulation
  - C++ custom build simulator
  - Goal: evaluate availability and associated revenue losses

  - We simulate failures and repairs of MEMS mirrors, splitters and WSSs based on component failure (repair) rates
  - Simulated time = 10<sup>9</sup> hours, encompassing over 5,000,000 simulated events
  - We consider two cases:
    - Shortest path (SP) routing algorithm when hard-wired nodes are used
    - Enforced Fibre Switching (EFS) routing algorithm when AoD

nodes are used.

![](_page_32_Picture_0.jpeg)

![](_page_32_Picture_1.jpeg)

![](_page_32_Picture_3.jpeg)

- Evaluation of end-to-end connections' successfulness for different total traffic loads.
- Using Monte Carlo simulation exponential distributed times to failure/repair of nodes and links are generated .
- When an AoD node component fails, node **spare bank** is checked for idle component identical to the failed type.
- If there is no idle component in the spare bank or in the case of a link failure, the recovery procedure is switched over to the **network level**, using **1+1 protection**.
- Simulated time = **10<sup>9</sup> hours**, with **10<sup>5</sup> simulated events**

# Results – Mean down time (MDT)

![](_page_33_Picture_1.jpeg)

![](_page_33_Figure_2.jpeg)

- The MDT obtained for following cases:
  - Hard-wired nodes with the shortest path (SP) and AoD nodes with SP/EFS
- Under low traffic AoD with SP achieves lower MDT
  - Some components remain idle used as redundancy
- Under high traffic– all components used
  - MDT of AoD with SP surpasses MDT of the hard-wired architecture
  - When SP is replaced with EFS, the MDT for AoD significantly decreases

![](_page_34_Picture_0.jpeg)

# Simulation results – HWnet vs. AoDnet

![](_page_34_Picture_2.jpeg)

At **low traffic** loads logical topologies are **not fully connected** higher no. of idle components.

At high traffic loads logical topologies are fully connected lower no. of idle components.

![](_page_34_Figure_5.jpeg)

HWnet FSP – HW network – with fixed shortest path routing

**AoDnet FSP** – AoD network – with fixed shortest path routing

![](_page_35_Picture_0.jpeg)

# Simulation results – Enforced fibre switching (EFS)

**DPP EFS** (dedicated path protection) + enforced fibre switching

EFS - Initial shortest path layout ⇒ transformed to optimised path routing in order to increase no. of fibreswitched (FS) paths ⇒ increase no. of redundant components.

![](_page_35_Figure_4.jpeg)

#### Trade off

Availability of prolonged paths vs 1 Availability of FS paths

![](_page_36_Picture_0.jpeg)

### Results – added redundant components

![](_page_36_Picture_2.jpeg)

![](_page_36_Figure_3.jpeg)

![](_page_37_Picture_0.jpeg)

![](_page_37_Picture_1.jpeg)

#### Results – GER topology

![](_page_37_Picture_3.jpeg)

- Average link distance 242 km
- MDT for AoD with DPP-OEFS is lower 25%

![](_page_37_Figure_6.jpeg)

![](_page_38_Picture_0.jpeg)

![](_page_38_Picture_2.jpeg)

- Average link distance 432 km
- In most cases, re-routed lightpaths can **NOT** compensate failure rate increase caused by extension of lightpath length.
- MDT approximately higher 30% over all test cases
  - **PROBLEM? HOW TO DECREASE MDT??**
  - DPP-FSP-RED added redundant WSS in 8-out-of-28 nodes
  - Cost increased 9%, while MDT was decreased 29%

![](_page_38_Figure_9.jpeg)

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![](_page_39_Picture_0.jpeg)

![](_page_39_Picture_1.jpeg)

![](_page_39_Picture_2.jpeg)

- The availability benefit of using AoD nodes with self-healing capability, compared to HW nodes, is evident at lower traffic loads using fixed shortest path (FSP) routing without additional investments in redundancy.
- If a tailored routing algorithm, with enforced fibre switching (EFS), is used, benefit of AoD deployment is extended to higher traffic loads.
- Always exists the possibility AoD self-healing capability can be improved by investement in redundant node components dedicated for failure recovery.